### Field Extensions and Elliptic Curves

Distributed Lab

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# Plan



- A bit of intuition
- General Definition
- Polynomial Fraction Rings
- Finite Field Extensions

### 2 Algebraic Closure

Definition

#### 3 Elliptic Curve

- Definition
- Group Structure

### Field Extensions

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 $\mathbb{Q}$  vs  $\mathbb{R}$ 

### Question #1

What is the difference between rational numbers  $\mathbb Q$  and real numbers  $\mathbb R$ ?

### Definition

**Rational numbers**  $\mathbb{Q}$  are defined as the set  $\{\frac{n}{m} : n \in \mathbb{Z}, m \in \mathbb{N}\}$ .

### Question #2

Why cannot we say  $m \in \mathbb{Z}$ , similarly to n?

### Theorem

 $\sqrt{2}$  is not a rational number. Neither is  $\pi$  and e. But they are reals.

### Conclusion

 ${\mathbb R}$  is sort of "an extended version of  ${\mathbb Q}$  ".

### **Rethorical Question**

Can we extend  $\mathbb{R}$ ?

Yes — just use complex numbers  $\mathbb{C}!$ 

### Definition

Complex numbers  $\mathbb{C}$  is defined as the set of x + iy where  $i^2 = -1$ .

### Definition

Complex numbers  $\mathbb{C}$  are defined as the set of pairs  $(x, y) \in \mathbb{R}^2$  where addition is defined as  $(x_1, y_1) + (x_2, y_2) = (x_1 + x_2, y_1 + y_2)$ , and the multiplication is:

$$(x_1, y_1) \cdot (x_2, y_2) = (x_1x_2 - y_1y_2, x_1y_2 + x_2y_1).$$

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# A bit about complex numbers

#### Theorem

 $(\mathbb{C},+,\times)$  is a field.

#### Example

Let us see how arithmetic is performed in  $\mathbb{C}$ .

- Addition: (2+3i) + (4+5i) = 6+8i.
- Multiplication:  $(1+i)(2+i) = 2 + 3i + i^2 = 1 + 3i$ .
- Division:

$$\frac{2+i}{1+i} = \frac{(2+i)(1-i)}{(1+i)(1-i)} = \frac{2-i-i^2}{1-i^2} = \frac{3-i}{2} = \frac{3}{2} - \frac{1}{2}i$$

### Question

What is (1 + i) + (2 + i)? i(1 + i)? 1/i?

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# Field Extension

### ${\sf Conclusion} + {\sf Question}$

 $\mathbb C$  is sort of "an extended version of  $\mathbb R".$  Thus, we have

 $\mathbb{Q} \subset \mathbb{R} \subset \mathbb{C}$ , where  $\mathbb{Q}, \mathbb{R}, \mathbb{C}$  are fields

So we have two questions in mind:

- Is there any mathematical term for this?
- Can we go further?

### Definition

Let  $\mathbb{F}$  be a field. A field  $\mathbb{K}$  is called an **extension** of  $\mathbb{F}$  if  $\mathbb{F} \subset \mathbb{K}$  which we denote as  $\mathbb{K}/\mathbb{F}$ .

#### Example

 $\mathbb{C}/\mathbb{R}$  is a field extension. So is  $\mathbb{R}/\mathbb{Q}$ .

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 $\mathbb{Q}(\sqrt{2})$  and  $\mathbb{Q}(i)$ 

#### Example

Define  $\mathbb{Q}(\sqrt{2}) := \{p + q\sqrt{2} : p, q \in \mathbb{Q}\}$ . This is a field extension of  $\mathbb{Q}$ . Arithmetic over  $\mathbb{Q}(\sqrt{2})$  looks like:

- Addition:  $(1+2\sqrt{2}) + (3+4\sqrt{2}) = 4 + 6\sqrt{2}$ .
- Multiplication:  $(1+2\sqrt{2})(1+\sqrt{2}) = 1+3\sqrt{2}+2\sqrt{2}^2 = 5+3\sqrt{2}$ .

Division:

$$\frac{1+2\sqrt{2}}{1+\sqrt{2}} = \frac{(1+2\sqrt{2})(\sqrt{2}-1)}{(\sqrt{2}+1)(\sqrt{2}-1)}$$

#### Example

Similarly,  $\mathbb{Q}(i) := \{p + qi : p, q \in \mathbb{Q}\}$  is a field extension of  $\mathbb{Q}$ .

# $\mathbb{Q}(\sqrt{2},i)$

#### Example

Define  $\mathbb{Q}(\sqrt{2}, i) = \{\alpha + \beta\sqrt{2} : \alpha, \beta \in \mathbb{Q}(i)\}$ . Typicall element of  $\mathbb{Q}(\sqrt{2}, i)$  can be written as:

$$(a + bi) + (c + di)\sqrt{2} = a + c\sqrt{2} + b\sqrt{2}i + di\sqrt{2}$$

#### Notice

Each element of  $\mathbb{Q}(\sqrt{2}, i)$  is a linear combination of  $\{1, \sqrt{2}, i, \sqrt{2}i\}$ . This is usually called a **basis**. Moreover, to denote the dimensionality of  $\mathbb{Q}(\sqrt{2}, i)$  over  $\mathbb{Q}$ , we write  $[\mathbb{Q}(\sqrt{2}, i) : \mathbb{Q}] = 4$ .

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### Definition... "Kinda"

Consider the set  $\mathcal{P}$  — a set of polynomials  $\mathbb{R}[x]$  modulo  $p(x) := x^2 + 1$ .

#### Example

For example, 1, 5 + x, 3x,  $1 + 2x \in \mathcal{P}$ .

But what about  $x^2 + 2x + 4$ ? We can divide by  $x^2 + 1$ !

$$x^{2} + 2x + 4 = (x^{2} + 1) \cdot 1 + (2x + 3)$$

So in P, we have  $x^2 + 2x + 4 = 2x + 3$ .

# Real Polynomials modulo $x^2 + 1$

#### Arithmetic

Over this field, we can do arithmetic as usual.

- Addition: (1 + x) + (2 + 3x) = 3 + 4x.
- Multiplication:  $(1 + x)(2 + x) = x^2 + 3x + 2 = 3x + 1$ .
- Inverse:

$$\left(\frac{1+x}{2}\right)^{-1} = 1 - x$$

Indeed,

$$\frac{1+x}{2} \cdot (1-x) = \frac{1}{2} \cdot (1-x^2) = \frac{1}{2} \left( -(x^2+1)+2 \right) = 1 \text{ (in } \mathcal{P})$$

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### Hold on a minute...

### Results

• 
$$(1+x) + (2+3x) = 3+4x$$

• 
$$(1+x)(2+x) = 1 + 3x$$

• 
$$\left(\frac{1+x}{2}\right)^{-1} = 1-x$$

### Same, but over $\ensuremath{\mathbb{C}}$

Let us do the same, but instead of X, use i.

• 
$$(1+i) + (2+3i) = 3+4i$$
.  
•  $(1+i)(2+i) = 2+3i+i^2 = 1+3i$   
•  $\frac{1}{\frac{1+i}{2}} = \frac{2}{1+i} = \frac{2(1-i)}{(1+i)(1-i)} = 1-i$ .

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### Hold on a minute...



So, basically,  $\mathcal{P}$  and  $\mathbb{C}$  have the same structure! Formally, they are isomorphic:  $\mathcal{P} \cong \mathbb{C}$ .

#### Question

Could we have used  $x^2 + 3$  instead of  $x^2 + 1$ ? What about  $x^2 + x + 1$ ?

Yes, any **irreducible** 2nd-degree polynomial p(x) over  $\mathbb{R}$  can be used. Typically, this is denoted as  $\mathbb{R}[x]/(p(x))$ .

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# Isomorphisms

### Reminder

For two groups  $(\mathbb{G}, +)$  and  $(\mathbb{H}, \times)$  we defined homomorphism to be a function  $\phi : \mathbb{G} \to \mathbb{H}$  such that

$$\phi(\mathbf{a} + \mathbf{b}) = \phi(\mathbf{a}) \times \phi(\mathbf{b})$$

However, we claim that  $\mathbb{R}/(x^2+1)\cong\mathbb{C}$ , which are fields, not groups.

### Definition

A field isomorphism is a function  $\phi : (\mathbb{F}, +, \times) \to (\mathbb{K}, \oplus, \otimes)$  such that

- $\phi(a+b) = \phi(a) \oplus \phi(b)$
- $\phi(a \times b) = \phi(a) \otimes \phi(b)$
- $\phi(1_{\mathbb{F}}) = 1_{\mathbb{K}}$

But for now, "congruence" essentially means "exhibit the same structure".

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#### Theorem

Let  $\mathbb{F}$  be a field and  $\mu(x)$  — irreducible polynomial over  $\mathbb{F}$  (reduction polynomial). Consider a set of polynomials over  $\mathbb{F}[x]$  modulo  $\mu(x) \in \mathbb{F}[x]$ , formally denoted as  $\mathbb{F}[x]/(\mu(x))$ . Then,  $\mathbb{F}[x]/(\mu(x))$  is a field.

#### Theorem

Let  $\mathbb{F}$  be a field and  $\mu \in \mathbb{F}[X]$  is an irreducible polynomial of degree n and let  $\mathbb{K} := \mathbb{F}[X]/(\mu(X))$ . Let  $\theta \in \mathbb{K}$  be the root of  $\mu$  over  $\mathbb{K}$ . Then,

$$\mathbb{K} = \{c_0 + c_1\theta + \cdots + c_{n-1}\theta^{n-1} : c_0, \ldots, c_{n-1} \in \mathbb{F}\}$$

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# Coming back to previous examples

#### Example

Again, consider 
$$\mathbb{Q}(\sqrt{2}) = \{q + p\sqrt{2} : p, q \in \mathbb{Q}\}$$
. Then,

$$\mathbb{Q}(\sqrt{2}) \cong \mathbb{Q}[x]/(x^2-2)$$

#### Example

Similarly,  $\mathbb{Q}(i) \cong \mathbb{Q}[x]/(x^2+1)$ .

#### Example

And  $\mathbb{Q}(\sqrt{2}, i)$  is just a little bit more tricky. Notice that we can take

$$p(x) := (x^2 - 2)(x^2 + 1) = x^4 - x^2 - 2$$

So  $\mathbb{Q}(\sqrt{2}, i) \cong \mathbb{Q}[x]/(x^4 - x^2 - 2).$ 

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# Finite Field Extension

#### Definition

Recall that  $\mathbb{F}_p$  (prime field) is a set  $\{0, 1, \dots, p-1\}$  with arithmetic modulo p.

In many cases, we need to extend  $\mathbb{F}_p$  2, 4, 8, 12, 24 times. For this, we use the so-called **finite field extension**.

### Definition

Suppose p is prime and  $m \ge 2$ . Let  $\mu \in \mathbb{F}_p[X]$  be an irreducible polynomial of degree m. Then, elements of  $\mathbb{F}_{p^m}$  are polynomials in  $\mathbb{F}_p^{(\le m)}[X]$  modulo  $\mu(x)$ . In other words,

$$\mathbb{F}_{p^m} = \{c_0 + c_1 X + \cdots + c_{m-1} X^{m-1} : c_0, \ldots, c_{m-1} \in \mathbb{F}_p\},\$$

where all operations are performed modulo  $\mu(X)$ .

# Examples

It would be convenient to build  $\mathbb{F}_{p^2}$  as  $\mathbb{F}_p[i]/(i^2+1)$ , but is it always possible? In other words, when  $X^2 = -1$  has a solution in  $\mathbb{F}_p$ ?

#### Theorem

Let p be an odd prime. Then  $X^2 + 1$  is irreducible in  $\mathbb{F}_p[X]$  if and only if  $p \equiv 3 \pmod{4}$ .

#### Example

Pick p = 19. Then  $\mathbb{F}_{361} := \mathbb{F}_{19}[i]/(i^2 + 1)$ . So typical elements are:

1+3i, 10+15i, 18+18i, 5, 7i, ...

• Addition: (1+10i) + (18+15i) = 19 + 25i = 6i.

#### • Multiplication:

 $(5+6i)(6+7i) = 30 + 71i + 42i^2 = -12 + 71i = 7 + 14i.$ 

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# More Examples: Binary Extension Fields

#### Example

Consider the  $\mathbb{F}_{2^4}.$  Then, there are 16 elements in this set:

$$\begin{array}{l} 0,1,X,X+1,\\ X^2,X^2+1,X^2+X,X^2+X+1,\\ X^3,X^3+1,X^3+X,X^3+X+1,\\ X^3+X^2,X^3+X^2+1,X^3+X^2+X,X^3+X^2+X+1. \end{array}$$

Set  $\mu(X) := X^4 + X + 1$ . Then, operations are performed in the following manner:

- Addition:  $(X^3 + X^2 + 1) + (X^2 + X + 1) = X^3 + X$ .
- Multiplication:  $(X^3 + X^2 + 1) \cdot (X^2 + X + 1) = X^2 + 1$  since:
- Inversion:  $(X^3 + X^2 + 1)^{-1} = X^2$  since  $(X^3 + X^2 + 1) \cdot X^2 \mod (X^4 + X + 1) = 1.$

## More Examples: BN254

#### Example

Consider the BN254 scalar field, used in SNARKs:

 $p = 0 \times 30644 \text{e}72 \text{e}131 \text{a} 029 \cdots \text{a}8 \text{d}3 \text{c}208 \text{c}16 \text{d}87 \text{c}\text{f} \text{d}47$ 

- Then,  $\mathbb{F}_{p^2} := \mathbb{F}_p[u]/(u^2+1)$  since  $p \equiv 3 \pmod{4}$ .
- Define  $\xi := 9 + u \in \mathbb{F}_{p^2}$ . Then, set  $\mathbb{F}_{p^6} := \mathbb{F}_{p^2}[v]/(v^3 \xi)$ .
- Finally, set  $\mathbb{F}_{p^{12}} := \mathbb{F}_{p^6}[w]/(w^2 v)$ .

Equivalently, we can write:

$$\mathbb{F}_{p^{12}} := \mathbb{F}_p[w] / (w^{12} - 18w^6 + 82)$$

### Algebraic Closure

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# Definition

### Definition

A field  $\mathbb{F}$  is called **algebraically closed** if every non-constant polynomial  $p(x) \in \mathbb{F}[X]$  has a root in  $\mathbb{F}$ .

#### Example

 $\mathbb{R}$  is not algebraically closed since  $X^2 + 1$  has no roots in  $\mathbb{R}$ . However,  $\mathbb{C}$  is algebraically closed, which follows from the fundamental theorem of algebra. Since  $\mathbb{C}$  is a field extension of  $\mathbb{R}$ , it is also an algebraic closure of  $\mathbb{R}$ . This is commonly denoted as  $\overline{\mathbb{R}} = \mathbb{C}$ .

#### Definition

A field  $\mathbb{K}$  is called an **algebraic closure** of  $\mathbb{F}$  if  $\mathbb{K}/\mathbb{F}$  is algebraically closed. This is denoted as  $\overline{\mathbb{F}} = \mathbb{K}$ .

# Algebraic Closure for Finite Fields

Recall that we are cryptographers, not mathematicials. So we are interested in  $\overline{\mathbb{F}}_p$ . So I have two news to you:

- Good news:  $\overline{\mathbb{F}}_p$  exists.
- Bad news:  $\overline{\mathbb{F}}_p$  is infinite.

#### Theorem

No finite field  $\mathbb{F}$  is algebraically closed.

**Proof.** Suppose  $f_1, f_2, \ldots, f_n \in \mathbb{F}$  are all elements of  $\mathbb{F}$ . Consider the following polynomial:

$$p(x) = \prod_{i=1}^{n} (x - f_i) + 1 = (x - f_1)(x - f_2) \cdots (x - f_n) + 1.$$

Clearly, p(x) is a non-constant polynomial and has no roots in  $\mathbb{F}$ , since for any  $f \in \mathbb{F}$ , one has p(f) = 1.

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So what?

But what form does the  $\overline{\mathbb{F}}_p$  have? Well, it is a union of all  $\mathbb{F}_{p^k}$  for  $k \ge 1$ . This is formally written as:

$$\overline{\mathbb{F}}_{p} = \bigcup_{k \in \mathbb{N}} \mathbb{F}_{p^{k}}$$

#### Remark

But this definition is super counter-intuitive! So here how we usually interpret it. Suppose I tell you that polynomial q(x) has a root in  $\overline{\mathbb{F}}_p$ . What that means is that there exists some extension  $\mathbb{F}_{p^m}$  such that for some  $\alpha \in \mathbb{F}_{p^m}$ ,  $q(\alpha) = 0$ . We do not know how large this *m* is, but we know that it exists. For that reason,  $\overline{\mathbb{F}}_p$  is defined as an infinite union of all possible field extensions.

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# Elliptic Curve

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# Definition

### Definition

Suppose that  $\mathbb{K}$  is a field. An **elliptic curve** *E* over  $\mathbb{K}$  is defined as a set of points  $(x, y) \in \mathbb{K}^2$ :

$$y^2 = x^3 + ax + b,$$

called a **Short Weierstrass equation**, where  $a, b \in \mathbb{K}$  and  $4a^3 + 27b^2 \neq 0$ . We denote  $E/\mathbb{K}$  to denote the elliptic curve over field  $\mathbb{K}$ .

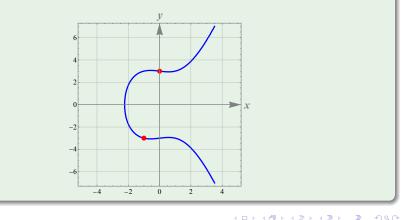
#### Definition

We say that  $P = (x_P, y_P) \in \mathbb{A}^2(\mathbb{K})$  is the **affine representation** of the point on the elliptic curve  $E/\mathbb{K}$  if it satisfies the equation  $y_P^2 = x_P^3 + ax_P + b$ .

### Examples

### Example

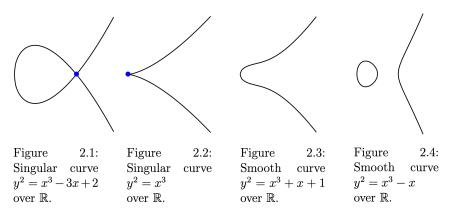
Consider  $E/\mathbb{Q}$ :  $y^2 = x^3 - x + 9$ . Valid affine points on  $E/\mathbb{Q}$  are, for example,  $P = (0,3), Q = (-1,-3) \in \mathbb{A}^2(\mathbb{Q})$ .



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# More Examples

Some more examples<sup>1</sup>:



<sup>1</sup> Figure taken from "Pair	ings for Beginners"	□ → 《圖 → 《圖 → 《圖 →	≣
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# Real Elliptic Curves

But real elliptic curves are not that simple. Here how they look like<sup>2</sup>:

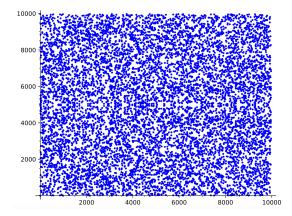


Figure: Curve  $E/\mathbb{F}_{9973}$ :  $y^2 = x^3 - 2x + 1$  over the finite field

<sup>2</sup> Figure taken from "M	loonmath"		₹ <b>୬</b> ९२
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# Defining a Group Structure: A Few Words

### Definition

The set of points on the curve, denoted as  $E_{a,b}(\mathbb{K})$ , is defined as:

$$E_{a,b}(\mathbb{K}) = \{(x,y) \in \mathbb{A}^2(\mathbb{K}) : y^2 = x^3 + ax + b\} \cup \{\mathcal{O}\},\$$

where  $\mathcal{O}$  is the so-called **point at infinity**.

#### Remark #1

If 
$$(x_P, y_P) \in E(\mathbb{K})$$
 then  $(x_P, -y_P) \in E(\mathbb{K})$ .

### Remark #2

Typically,  $\mathbb{K} = \overline{\mathbb{F}}_p$ : we do not conretize over which finite field we define the elliptic curve.

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### Defining a Group Structure: Chord Method

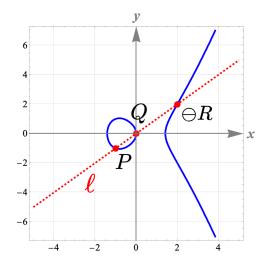


Figure: Chord method for adding two points

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### Defining a Group Structure: Tangent Method

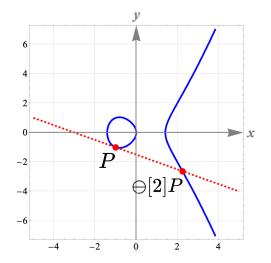


Figure: Tangent method for the point doubling

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### Idea of Derivation

Line equation through  $P = (x_P, y_P), Q = (x_Q, y_Q)$ :

$$\ell: y = \lambda(x - x_P) + y_P, \ \lambda = \frac{y_Q - y_P}{x_Q - x_P}$$

So all we need is to solve the system of equations:

$$\begin{cases} y^2 = x^3 + ax + b\\ y = \lambda(x - x_P) + y_P \end{cases}$$

Substituting y from the second equation to the first one, we get a cubic equation. Using Vieta's formula, one can derive

$$x_P + x_Q + x_R = \lambda^2$$

The rest is easy to finish.

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# Group Law

#### Definition

- **1** Point at infinity  $\mathcal{O}$  is an identity element.
- **3** If  $x_P \neq x_Q$ , use the **chord method**. Define  $\lambda := \frac{y_P y_Q}{x_P x_Q}$  the slope between *P* and *Q*. Set the resultant coordinates as:

$$x_R := \lambda^2 - x_P - x_Q, \quad y_R := \lambda(x_P - x_R) - y_P.$$

3 If  $x_P = x_Q$  and  $y_P = y_Q$  (that is, P = Q), use the **tangent method**. Define the slope of the tangent at P as  $\lambda := \frac{3x_P^2 + a}{2y_P}$  and set

$$x_R := \lambda^2 - 2x_P, \quad y_R := \lambda(x_P - x_R) - y_P.$$

• Otherwise, define  $P \oplus Q := O$ .

# One more Illustration

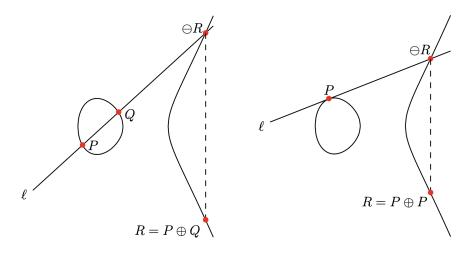


Figure 2.5: Elliptic curve addition.

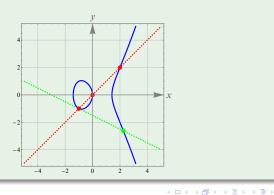
Figure 2.6: Elliptic curve doubling.

# Example

#### Example

Consider  $E/\mathbb{R}$ :  $y^2 = x^3 - 2x$ .

- Addition:  $(-1,1) \oplus (0,0) = (2,-2), (2,2) \oplus (-1,-1) = (0,0).$
- Doubling:  $[2](-1,-1) = (\frac{9}{4},-\frac{21}{8}).$



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#### Theorem

 $(E(\overline{\mathbb{F}}),\oplus)$  forms an abelian group.

Now, let us consider the group order  $r := |E(\mathbb{F}_{p^m})|$ .

### Theorem

Hasse's Theorem on Elliptic Curves.  $r = p^m + 1 - t$  for some integer  $|t| \le 2\sqrt{p^m}$ . A bit more intuitive explanation: the number of points on the curve is close to  $p^m + 1$ . The value t is called the trace of Frobenius.

#### Remark

In fact,  $r = |E(\mathbb{F}_{p^m})|$  can be computed in  $O(\log(p^m))$ , so the number of points can be computed efficiently even for fairly large primes p.

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# Discrete Logarithm

#### Definition

Let  $P \in E(\overline{\mathbb{F}}_p)$  and  $\alpha \in \mathbb{Z}_r$ . Define the scalar multiplication  $[\alpha]P$  as adding P to itself  $\alpha - 1$  times (also set  $[0]P := \mathcal{O}$ ).

#### Definition

Suppose *E* is cyclic, meaning,  $\langle G \rangle = E$  for some  $G \in E$ . The **discrete logarithm problem** on *E* consists in the following: suppose  $P = [\alpha]G$  for some  $\alpha \in \mathbb{Z}_r$ . Find  $\alpha$  based on *P*.

#### Remark

If r is a product of primes  $p_1, p_2, \ldots, p_k$  such that  $p_1 < p_2 < \cdots < p_k$ , then the best-known algorithm to solve the discrete logarithm problem is no significantly better than  $O(\sqrt{p_1})$ .

August 1, 2024

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### Thank you for your attention!

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Image: A matrix